

Fundamentals of Modern Cryptography

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This presentation explains the fundamentals of modern cryptographic methods.

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What Cryptography Can Accomplish

Authenticity Verify who is on the other end of the communication channel

Confidentiality Only the other party can read the original messages

Integrity No other party can change or add messages

Other features are often desired, e.g., non-repudiation.

Outline

1. Ciphers and hashes
2. Big numbers
3. Primes
4. Private communication
5. Public/private-key communication
6. Conclusion

1 Ciphers and Hashes: Ciphers

encrypt(message, secret) = cipher

decrypt(cipher, secret) = message

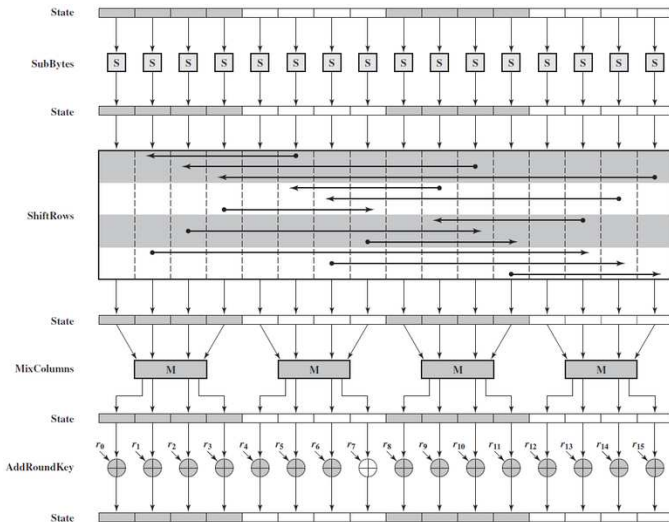
decrypt(encrypt(message, secret), secret) = message

Ciphers map a message to a same-length unique ciphertext, which can be reversed.

AES (Advanced Encryption Standard)

- ▶ 128-bit block cipher
- ▶ Supports key lengths of 128, 192, and 256 bits, e.g., AES128 uses a 128-bit key
- ▶ Uses confusion (substitution) and diffusion; see https://en.wikipedia.org/wiki/Confusion_and_diffusion
- ▶ No known attack except for exhaustive key search (this is ideal)
- ▶ AES-specific CPU instructions speed processing 3-10x, e.g., AES-NI (check `/proc/cpuinfo` for “aes”)
- ▶ Elegant algorithm based on polynomials over finite fields, see https://en.wikipedia.org/wiki/Advanced_Encryption_Standard
- ▶ Trusted by the US government for top secret communication

One Cycle of AES



AES128 requires 10 cycles.

https://commons.wikimedia.org/wiki/File:AES_Encryption_Round.png

Cipher Modes

message = 128bit₁ || 128bit₂ || ...

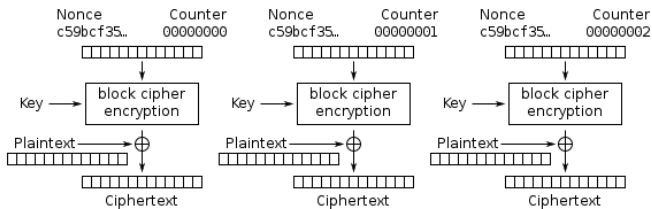
encrypt(message, secret) = *encrypt(128bit₁, secret)* || *encrypt(128bit₂, secret)* || ...

block encrypt(128bit₁, secret) = *cipher(secret, 128bit₁ XOR (previous output OR nonce))* (CBC)

stream encrypt(128bit₁, secret) = 128bit₁ XOR *cipher(secret, counter || nonce)* (CTR, GCM)

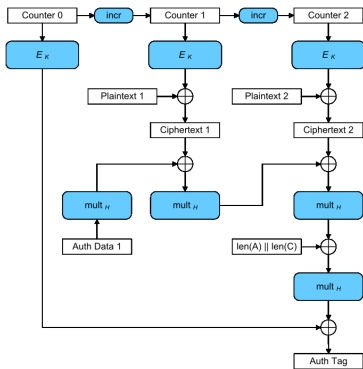
Stream encryption built on block ciphers is particularly useful for streaming protocols, e.g., SSH. You can generate a 128-bit encryption block and XOR the 16 bytes individually.

CTR (Counter) Illustrated



https://commons.wikimedia.org/wiki/File:CBC_encryption.svg

GCM (Galois Counter Mode) Illustrated



GCM is similar to CTR mode but also creates an integrity tag that can be used to verify that the entire message was created with the same secret key and unmodified.

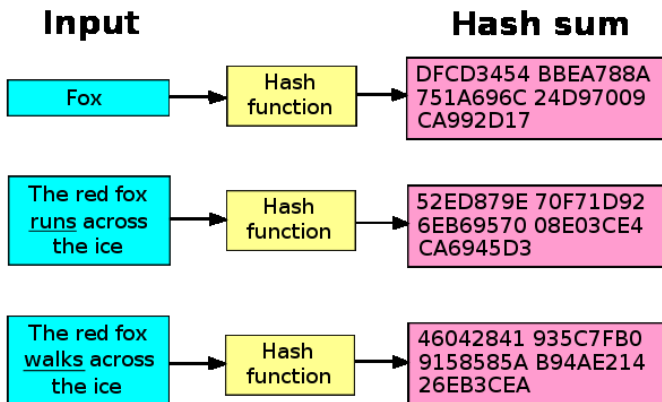
https://en.wikipedia.org/wiki/Galois/Counter_Mode#/media/File:GCM-Galois_Counter_Mode.svg

Hash

$$\begin{aligned} \mathit{hash}(\mathit{message}) &= \mathit{hash} \\ \mathit{function}(\mathit{hash}) &\neq \mathit{message} \end{aligned}$$

A hash maps a variable-length value to a fixed-length value, with minimal collisions. Collisions are likely at $\sqrt{\mathit{hash\ space}}$, e.g., a hash algorithm with 2^{256} possible outputs is likely to generate a collision among 2^{128} outputs. The hash name indicates the hash output bit length, e.g., SHA-256, SHA-384.

Hashing Illustrated



https://commons.wikimedia.org/wiki/File:Hash_function_long.svg

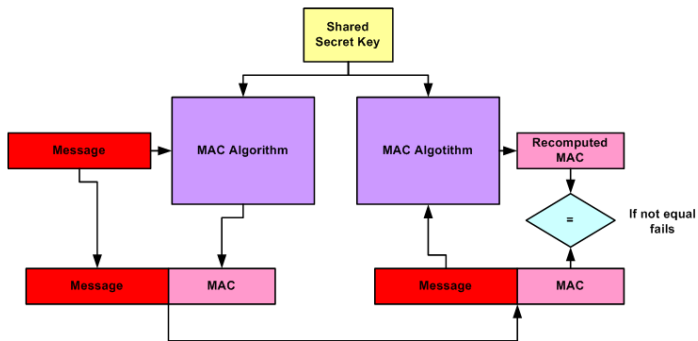
MAC (Message Authentication Code)

$hash(message, secret) = hash$

$hash(message, no\ secret) \neq matching\ hash$

- ▶ Hashing (or limited block encryption) allows you to safely prove you know a secret
- ▶ Someone who also knows the secret can compute the same MAC and check that it matches
- ▶ Others who see the MAC can't reverse it to divulge the secret
- ▶ A message authentication code (MAC) proves that the message was created by someone who knows the secret
- ▶ Multiple pieces of information can be combined in a single hash, e.g.,
 - ▶ often the message, message length, and secret key are combined to produce a single hash value
 - ▶ this proves the message was produced by someone who knows the secret

MAC Illustrated



https://commons.wikimedia.org/wiki/File:Cryptographic_MAC_based_message_authentication.png

Split Keys

While it is possible to merge secrets in a single hash, it is also possible to split a secret into parts so all parts are needed to reconstruct the secret, e.g two parts:

$$\begin{aligned} \textit{key1} &= \textit{random number} \\ \textit{master secret XOR secret1} &= \textit{secret2} \\ &\textit{to reconstruct} \\ \textit{secret1 XOR secret2} &= \textit{master secret} \end{aligned}$$

This can be repeated to split a key into any number of parts; see https://en.wikipedia.org/wiki/Secret_sharing.

2 Big Numbers: Exponentiation

$$a^b = a \times a \times a \dots (b \text{ times})$$

e.g.

$$2^5 = 2 \times 2 \times 2 \times 2 \times 2$$

$$7\text{E}9 = 7 \times 10^9 = 7,000,000,000$$

Logarithms

$$a^b = c$$

$$b \times \log a = \log c$$

$$b = \frac{\log c}{\log a}$$

CPU Register Sizes

$$32 \text{ bit } 2^{32} \cong 4\text{E}9$$

$$64 \text{ bit } 2^{64} \cong 2\text{E}19$$

Real World Sizes

- ▶ Seconds in a billion years: 3×10^{16}
- ▶ Atoms in the universe: 1×10^{80}
- ▶ Possible chess games: 1×10^{120}

Cryptography Sizes

- ▶ AES128 (128 bit): $3E38$
- ▶ ECDHE prime (224 bit): $2E67$
- ▶ AES256, SHA256 (256 bit): $1E77$
- ▶ RSA prime (1024 bit): $2E308$
- ▶ DH prime, RSA composite (2048 bit): $4E616$
 - ▶ strength equivalent to a 112-bit symmetric cipher ($3E35$)
- ▶ Random values with range 2^b likely repeat after $2^{b/2}$ values
- ▶ Adding a bit doubles the strength, adding two bits quadruples the strength
- ▶ Doubling the number of bits exponentially increases the strength

This video explains the difficulty of breaking large computed secrets, https://www.youtube.com/watch?v=S9JGmA5_unY.

https://en.wikipedia.org/wiki/Key_size

3 Primes: Modulus

$$9 \text{ mod } 7 = 2$$

$$9 \% 7 = 2$$

Random number generation, prime number generation, and prime number detection are also required for cryptography.

Generator

6 raised to an integer power modulus 7 generates a subset of values less than 7:

$$6^0 \bmod 7 = 1$$

$$6^1 \bmod 7 = 6$$

$$6^2 \bmod 7 = 1$$

$$6^3 \bmod 7 = 6$$

Primitive Element

3 is a primitive element of 7 because it generates the full set:

$$3^1 \bmod 7 = 3$$

$$3^2 \bmod 7 = 2$$

$$3^3 \bmod 7 = 6$$

$$3^4 \bmod 7 = 4$$

$$3^5 \bmod 7 = 5$$

$$3^6 \bmod 7 = 1$$

$$3^7 \bmod 7 = 3$$

$$3^8 \bmod 7 = 2$$

Reverse

True for any integer n :

$$1 = 3^{6n} \pmod{7}$$

$$2 = 3^{6n+2} \pmod{7}$$

$$3 = 3^{6n+1} \pmod{7}$$

$$4 = 3^{6n+4} \pmod{7}$$

$$5 = 3^{6n+5} \pmod{7}$$

$$6 = 3^{6n+3} \pmod{7}$$

Exponentiation Commutativity

$$g^{x^y} = g^{x \times y} = g^{y \times x} = g^{y^x}$$

4 Private Communication: Ephemeral Diffie–Hellman (DHE)

Assume g is a (sub-group) generator of (safe) prime p :

| Alice | Bob |
|--|--|
| agree on g and p | agree on g and p |
| generate random x | generate random y |
| compute $g^x \bmod p$ | compute $g^y \bmod p$ |
| send $g^x \bmod p \rightarrow$ | \leftarrow send $g^y \bmod p$ |
| receive $g^y \bmod p \leftarrow$ | \rightarrow receive $g^x \bmod p$ |
| use x to compute $(g^y \bmod p)^x \bmod p$ | use y to compute $(g^x \bmod p)^y \bmod p$ |
| use as shared key $g^{xy} \bmod p$ | use as shared key $g^{xy} \bmod p$ |

Viewing transmission of g , p , $g^x \bmod p$, and $g^y \bmod p$ does not allow easy discovery of x , y , or $g^{xy} \bmod p$ (the secret key). This seems trivial until you realize it is being done using 2048-bit (3E616) values.

Exponentiation by high powers can be performed efficiently using **exponentiation by squaring** (binary exponentiation), optimized via **modular exponentiation**.

The Result

Both users have the same secret key:

$$g^{y^x} \bmod p = g^{x^y} \bmod p$$

while anyone viewing the exchange cannot derive the secret key. There is no information in the key. Instead, the computed key is used to generate a session key for encryption.

Fake Addition Key Exchange Example

Pretend you can't reverse the modulo computation (this is the effect of a large exponent):

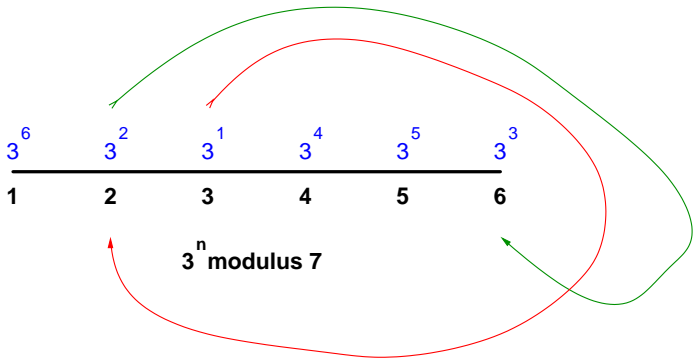
| Alice | Bob |
|---|---|
| agree on $g=6$ and $p=7$ generate random $x=5$ compute $(6+5) \bmod 7 = 4$ send $4 \rightarrow$ receive $3 \leftarrow$ use 3 to compute $(4 + 3) \bmod 7 = 0$ use as shared key 0 | agree on $g=6$ and $p=7$ generate random $y=4$ compute $(6+4) \bmod 7 = 3$ \leftarrow send 3 \rightarrow receive 4 use 4 to compute $(3 + 4) \bmod 7 = 0$ use as shared key 0 |

Real Key Exchange Example

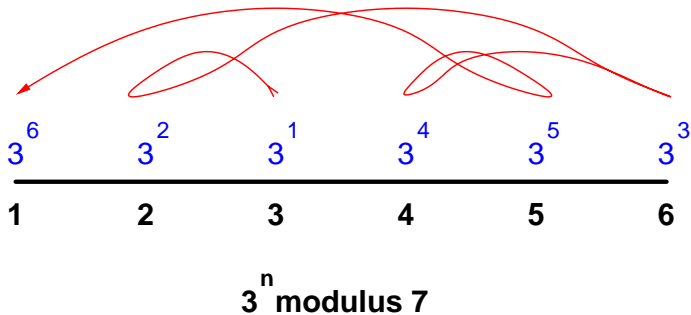
This uses very small values:

| Alice | Bob |
|--------------------------------------|--------------------------------------|
| agree on $g=3$ and $p=7$ | agree on $g=3$ and $p=7$ |
| generate random $x=3$ | generate random $y=4$ |
| compute $3^3 \bmod 7 = 6$ | compute $3^4 \bmod 7 = 4$ |
| send $6 \rightarrow$ | \leftarrow send 4 |
| receive $4 \leftarrow$ | \rightarrow receive 6 |
| use 4 to compute $6^4 \bmod 7 = 1$ | use 6 to compute $4^6 \bmod p = 1$ |
| use as shared key 1 | use as shared key 1 |

Unpredictability of Modulus Exponentiation



Unpredictability of Modulus Exponentiation



While it is computationally easy to compute higher exponents from a given modulus, it is computationally impossible to compute lesser exponent moduli (a trapdoor function).

Discrete Logarithm in a Finite Field

If an eavesdropper sees $g^x \bmod p$, how do they find x , e.g.,

$$3^x = 729$$

$$x \times \log 3 = \log 729$$

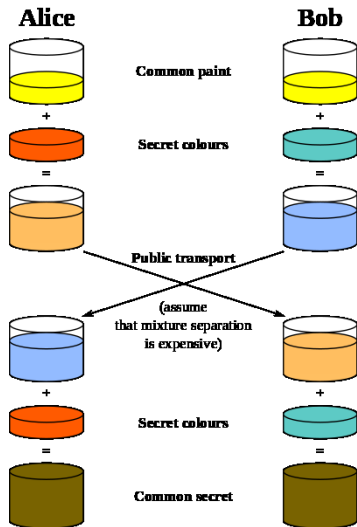
$$x = \frac{\log 729}{\log 3}$$

$$x = 6$$

$$(x \times \log 3) \bmod p \neq (\log 729) \bmod p$$

Modulus a prime p creates a finite (Galois) field where logarithms don't work; see https://en.wikipedia.org/wiki/Finite_field.

Diffie-Hellman Illustrated with Paint



https://commons.wikimedia.org/wiki/File:Diffie-Hellman_Key_Exchange.svg

Elliptic Curve Diffie–Hellman (ECDHE)

Similar to normal Diffie-Hellman because:

- ▶ A generator g and modulus-prime p are used, and are public
- ▶ Computation commutativity is used by each participant to privately compute the secret key

but it is different because:

- ▶ Values exist as points on an elliptic curve, rather than being linear (on a number line)
- ▶ g is an ordered-pair of integers (a 2D point), not a single integer

$$y^2 = x^3 + ax + b$$

Public constants a and b define the elliptic curve. Elliptic curves are Abelian groups, allowing commutativity; see https://en.wikipedia.org/wiki/Abelian_group.

ECDHE Exchange

- ▶ The participants each generate a random number, like DHE
- ▶ Starting from the generator point, they use point doubling and addition to simulate moving the point their random number of times
- ▶ They transmit their new points to each other
- ▶ Using the received points they double/add them their random number of times
- ▶ Both participants end at the same point on the elliptical curve
- ▶ They use the x component of the result to derive a shared secret

Elliptic Curve Doubling and Addition

Suppose the random number is 71 (binary 1000111). We need to compute the effect of starting at point g and advancing 70 more times along the elliptic curve. Doubling and **addition** can be used to shorten the computation to reach the final point:

$$2 \times g = 2g$$

$$2 \times 2g = 4g$$

$$2 \times 4g = 8g$$

$$2 \times 8g = 16g$$

$$16g + g = 17g$$

$$2 \times 17g = 34g$$

$$34g + g = 35g$$

$$2 \times 35g = 70g$$

$$70g + g = 71g$$

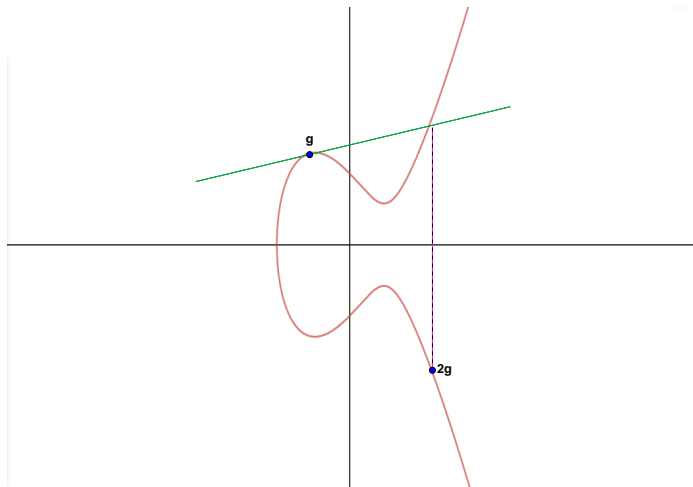
This is similar to *exponentiation by squaring* (https://en.wikipedia.org/wiki/Exponentiation_by_squaring), used by DHE

ECDHE Point Computations

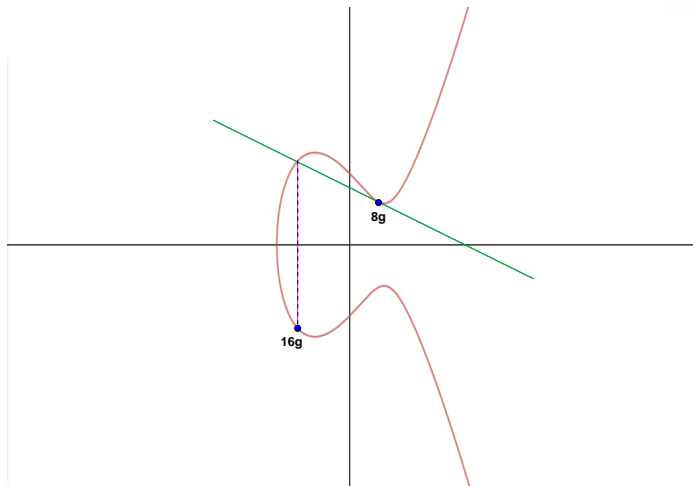
- ▶ Start with the generator point g
- ▶ Double it to get $2g$
- ▶ Doubling involves computing a line tangent to the point being doubled and finding the line intersection point with the elliptical curve
- ▶ This is effectively adding a point to itself, assuming the two points are on top of each other
- ▶ Reflect the point across the x axis
- ▶ Addition involves drawing a line between two points and computing the intersection point, then reflecting

<http://andrea.corbellini.name/2015/05/17/elliptic-curve-cryptography-a-gentle-introduction/>

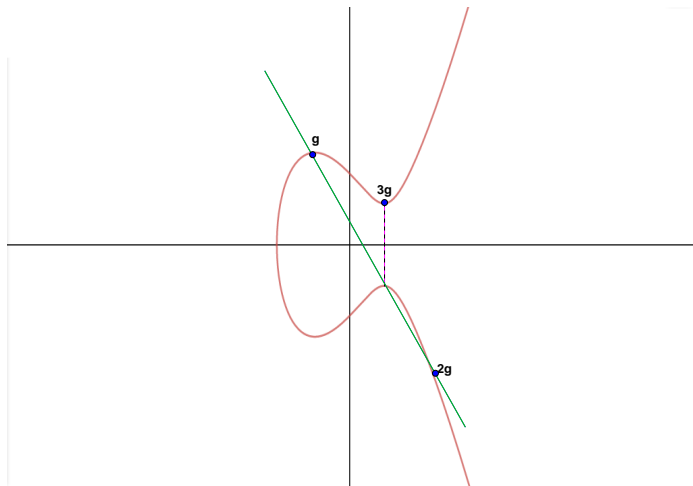
Elliptic Curve Doubling g



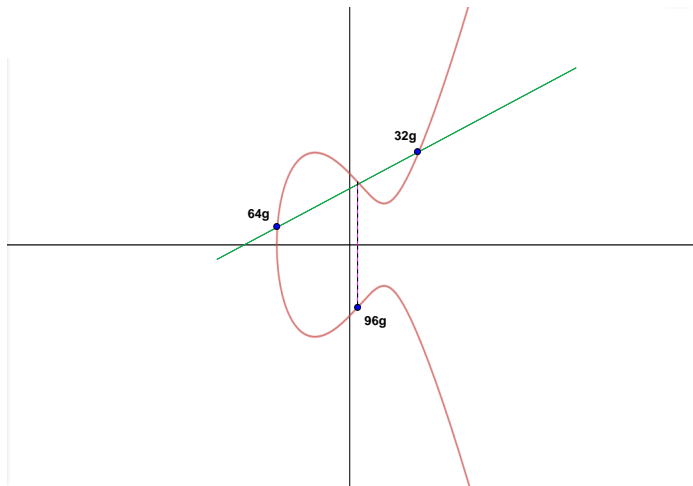
Elliptic Curve Doubling $8g$



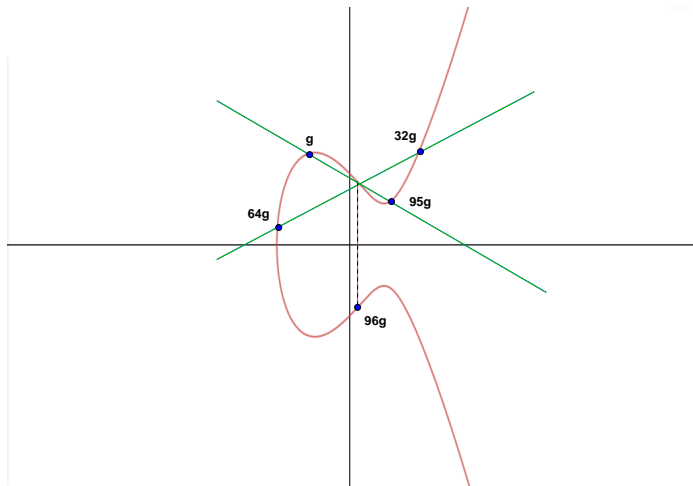
Elliptic Curve Adding g and $2g$



Elliptic Curve Adding $64g$ and $32g$

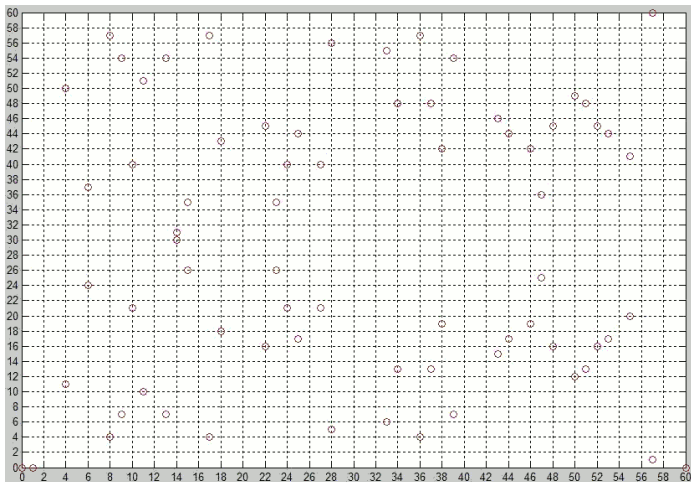


Adding One vs. Adding Larger Numbers



Elliptic Curve Distribution for Prime of 61

an



Modulus forces the values into a finite field.

https://hu.wikipedia.org/wiki/Elliptikus_g%C3%B6rbe#/media/File:Elliptic_curve_y%5E2%3Dx%5E3-x_on_finite_field_Z_61.PNG

ECDHE Conclusion

Elliptical curve Diffie-Hellman requires less computation than DHE, e.g., 224-bit ECDHE is as strong as 2048-bit DHE.

For an elliptical curve cryptography overview watch the first half of the video by Martijn Grooten, <https://youtu.be/yBr3Q6xiTw4>, and for details see the video by Nigel Smart, <https://youtu.be/t3JzdKE-Fhs>.

5 Public/Private-Key Communication: RSA (Rivest, Shamir, Adleman)

Diffie–Hellman:

- ▶ Negotiates a shared secret
- ▶ Uses a public exponent base g
- ▶ Uses a public modulus prime p
- ▶ Each client generates a secret exponent; is peer-to-peer key negotiation

RSA:

- ▶ Communicates a message
- ▶ The message is the exponent base
- ▶ The public modulus n is the product of two private primes, p and q
- ▶ One exponent e is public, another d is private
- ▶ Raising a message to the power e and then d (or the reverse order) returns the original message
- ▶ One party controls the private key
- ▶ Performs encryption and signing

Using Exponentiation to Return the Base

Fermat's Little Theorem, year 1640, states that for all m in the range $1 < m < p$ and p prime:

$$\begin{aligned}m^{p-1} \bmod p &= 1 \\m \times (m^{p-1} \bmod p) &= m \\(m \times m^{p-1}) \bmod p &= m \\m^p \bmod p &= m\end{aligned}$$

This allows large exponentiation in a modulus field to return the base value.

- ▶ <https://mathlesstraveled.com/2017/10/14/four-formats-for-fermat/>
- ▶ <https://mathlesstraveled.com/2017/11/13/fermats-little-theorem-proof-by-modular-arithmetic/>
- ▶ https://en.wikipedia.org/wiki/Proofs_of_Fermat%27s_little_theorem

Multiple Exponentiations Return the Base

Exponentiation in a modulus field is cyclic, so multiple exponent values return the same result. For integer k :

$$\begin{aligned}m^{p-1} \bmod p &= 1 \\m^{k \times (p-1)} \bmod p &= 1 \\m^{k \times (p-1) + 1} \bmod p &= m\end{aligned}$$

Using Two Primes to Return the Base

Euler's Theorem, year 1763, states that for primes p and q , and $lcm()$ as the least common multiple:

$$\begin{aligned}m^{p-1} \bmod p &= 1 \\m^{(p-1)(q-1)} \bmod (p \times q) &= 1 \\m^{lcm((p-1)(q-1))} \bmod (p \times q) &= 1 \\m^{lcm((p-1)(q-1))+1} \bmod (p \times q) &= m\end{aligned}$$

- ▶ https://en.wikipedia.org/wiki/Euler%27s_theorem
- ▶ https://en.wikipedia.org/wiki/Euler%27s_totient_function

Using Two Primes and Cyclic Exponentiation

For integer k :

$$\begin{aligned} m^{\text{lcm}((p-1)(q-1))+1} \bmod (p \times q) &= m \\ m^{k \times \text{lcm}((p-1)(q-1))+1} \bmod (p \times q) &= m \end{aligned}$$

Creating a Public/Private Exponent Pair

Instead of using p and q directly to compute an exponent pair that will return m , choose a small exponent e and compute an exponent d that will return m :

$$\begin{aligned}e \times d &= k \times \text{lcm}((p - 1)(q - 1)) + 1 \\(e \times d) \bmod \text{lcm}((p - 1)(q - 1)) &= 1 \\d &= \frac{1}{e} \bmod \text{lcm}((p - 1)(q - 1))*\end{aligned}$$

* Computing a fraction in a modulus field is accomplished using the Chinese remainder theorem,

https://en.wikipedia.org/wiki/Chinese_remainder_theorem.

Exponentiation Using d and e

$$\begin{aligned}m^{k \times \text{lcm}((p-1)(q-1))+1} \bmod (p \times q) &= m \\d \times e &= k \times \text{lcm}((p-1)(q-1)) + 1 \\m^{e \times d} \bmod (p \times q) &= m \\m^{e^d} \bmod (p \times q) &= m \\m^{d^e} \bmod (p \times q) &= m\end{aligned}$$

Intermediate Modulus Operations

$$m^{e^d} \bmod (p \times q) = m$$

$$m^{d^e} \bmod (p \times q) = m$$

$$(m^e \bmod (p \times q))^d \bmod (p \times q) = m$$

$$(m^d \bmod (p \times q))^e \bmod (p \times q) = m$$

Summary of Exponentiation With Modulus

- ▶ Base values raised to specific exponents with modulus return the original values, $m^p \bmod p = m$
- ▶ Multiple exponents can do this, $m^{k \times (p-1) + 1} \bmod p = m$
- ▶ When using a modulus that is the product of two primes, finding an exponent that returns the base value is hard to compute using just the modulus, $m^{\text{lcm}((p-1)(q-1)) + 1} \bmod (p \times q) = m$
- ▶ Multiple exponents can do this, $m^{k \times \text{lcm}((p-1)(q-1)) + 1} \bmod (p \times q) = m$
- ▶ Other exponent pairs, which are more distantly related to the two primes, can also do this, $e \times d = k \times \text{lcm}((p-1)(q-1)) + 1$
- ▶ Exponent pairs can be applied in any order,
 $m^{e \cdot d} \bmod (p \times q) = m^{d \cdot e} \bmod (p \times q)$
- ▶ Modulus can be applied to intermediate results,
 $(m^e \bmod (p \times q))^d \bmod (p \times q) = m$

Public/Private Designations

Public

- ▶ $n = p \times q$, used as a modulus
- ▶ exponent e

Private

- ▶ exponent d
- ▶ primes p and q
- ▶ $(p - 1)(q - 1)$, used to compute d

The RSA key length indicates the bits in n , e.g., 2048 bits.
Modulus is used to compute the exponent d and applied to the exponentiation result.

RSA Application

Dual-exponent ordering allows for encryption and signing (authentication):

$$\text{encrypt } message^e \bmod n = \text{encrypted}$$

$$\text{decrypt } encrypted^d \bmod n = \text{message}$$

$$\text{sign } hash(message)^d \bmod n = \text{signature}$$

$$\text{verify } signature^e \bmod n = hash(message)$$

Encryption uses the public exponent e to encrypt a message to produce encrypted text that is later decrypted with private exponent d . Signing uses the private exponent d to encrypt the hash of a message to produce a signature which is later validated by applying the public exponent e . Signing the hash of Diffie–Hellman parameters authenticates the parameters as being created by the private key owner.

6. Conclusion: What Cryptography Can Accomplish

Authenticity Verify who is on the other end of the communication channel (RSA via X.509 certificates)

Confidentiality Only the other party can read the original messages (ECDHE, AES)

Integrity No other party can change or add messages (GCM, SHA)

Further Cryptography Study

For additional study about cryptography, watch the nine video lessons at <https://www.youtube.com/playlist?list=PLqhpVxkBo1dPiKHym2Cx0KEnqC0350JH2>. The videos were created by Bill Buchanan, Professor in the School of Computing at Edinburgh Napier University. His website, <http://asecuritysite.com/>, contains a wealth of information about digital security.

Further Cryptography Reading

- ▶ Niels Ferguson, Bruce Schneier, Tadayoshi Kohno, *Cryptography Engineering: Design Principles and Practical Applications*, 2010: practical overview of cryptographic protocols
- ▶ Christof Paar, Jan Pelzl, *Understanding Cryptography: A Textbook for Students and Practitioners*, 2010: scholarly overview of all modern cryptographic protocols
- ▶ Bruce Schneier, *Applied Cryptography*, 1996: historical but broad overview of the cryptographic landscape

Conclusion



<http://momjian.us/presentations>

<https://www.flickr.com/photos/vmax137/>